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# The Trimmed Iterative Closest Point Algorithm

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# Euclidean alignment of two 3D point sets

- Given two roughly pre-registered 3D point sets,  $\mathcal{P}$  (data) and  $\mathcal{M}$  (model), find shift & rotation that bring  $\mathcal{P}$  into **best possible alignment** with  $\mathcal{M}$ .  $\triangleright\triangleright$
- **Applications:**
  - 3D model acquisition – reverse engineering, scene reconstruction
  - motion analysis – model-based tracking
- **Problems:**
  - partially overlapping point sets – incomplete measurements
  - noisy measurements
  - erroneous measurements – **outliers**
  - shape defects

Note: We consider **Euclidean** alignment. Other alignments, e.g., **affine**, are also studied.

# Iterative Closest Point (ICP) algorithm

Besl and McKay (1992): Standard solution to alignment problem.

## *Algorithm 1: Iterative Closest Point*

1. **Pair each point** of  $\mathcal{P}$  to closest point in  $\mathcal{M}$ .
  2. **Compute motion** that minimises mean square error (MSE) between paired points.
  3. **Apply motion** to  $\mathcal{P}$  and update MSE.
  4. **Iterate** until convergence.
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Chen and Medioni (1992): A similar iterative scheme using different pairing procedure based on surface normal vector.

**We use Besl's formulation:** applicable to volumetric measurements.

# Properties of ICP

- **Pre-registration** required:
  - manual
  - known sensor motion between two measurements
- **Point pairing**:
  - computationally demanding
  - special data structures used to speed up (k-D trees, spatial bins)
- **Optimal motion**: closed-form solutions available
- Proved to **converge** to a local minimum
- Applicable to **surface** as well as **volumetric** measurements
- Drawbacks:
  - **not robust**: assumes outlier-free data and  $\mathcal{P} \subset \mathcal{M}$
  - converges quite slowly

# Closed-form solutions for optimal rigid motion

- Unit Quaternions (Horn 1987): used by original ICP and our TrICP
- Singular Value Decomposition (Arun 1987)
- Orthogonal Matrices (Horn 1988)
- Dual Quaternions (Walker 1991)

*Properties of the methods (Eggert 1997)*

Method	Accuracy	2D Stability*	Speed, small $N_p$	Speed, large $N_p$
UQ	good	good	fair	fair
SVD	good	good	fair	fair
OM	fair	poor	good	poor
DQ	fair	fair	poor	good

\* Stability in presence of degenerate (2D) data.

# Existing variants of ICP

**Goal:** Improve robustness and convergence (speed).

**Categorisation criteria** (Rusinkiewicz 2001): How variants

1. **Select** subsets of  $\mathcal{P}$  and  $\mathcal{M}$ 
  - random sampling for a Monte-Carlo technique
2. **Match** (pair) selected points
  - closest point
  - in direction of normal vector: faster convergence when normals are precise
3. **Weight** and **reject** pairs
  - distribution of distances between paired points
  - geometric constraints (e.g., compatibility of normal vectors)
4. Assign **error metric** and **minimise** it
  - iterative: original ICP
  - direct: Levenberg-Marquardt algorithm

## Robustness and convergence: critical issues

ICP assumes that each point of  $\mathcal{P}$  has valid correspondence in  $\mathcal{M}$ .  
Not applicable to partially overlapping sets or sets containing outliers.

Previous attempts to robustify ICP: Reject wrong pairs based on

- **Statistical criteria.** Monte-Carlo type technique with robust statistics:
  - Least **median of squares** (LMedS)
  - Least **trimmed squares** (LTS)
- **Geometric criteria.** For example, Iterative Closest *Reciprocal* Point (Pajdla 1995) uses  $\epsilon$ -reciprocal correspondence:
  - **if** point  $\mathbf{p} \in \mathcal{P}$  has closest point  $\mathbf{m} \in \mathcal{M}$ , **then**
  - back-project  $\mathbf{m}$  onto  $\mathcal{P}$  by finding closest point  $\mathbf{p}' \in \mathcal{P}$
  - reject pair  $(\mathbf{p}, \mathbf{m})$  if  $\|\mathbf{p} - \mathbf{p}'\| > \epsilon$

Heterogeneous algorithms: heuristics combined, **convergence cannot be proved.**



# Robust statistics: LMedS and LTS

Sort distances between paired points, minimise

- **LMedS**: value in the middle of sorted sequence
  - operations **incompatible** with computation of optimal motion
- **LTS**: sum of certain number of least values (e.g., least 50%)
  - operations **compatible** with computation of optimal motion
  - better convergence rate, smoother objective function

Previous use of LMedS and LTS: Randomised robust regression

- estimate optimal motion parameters by repeatedly drawing random samples
- detect and reject outliers, find least squares solution for inliers

Robust to outliers, but breakdown point 50%  $\Rightarrow$  **minimum overlap** 50%.

# Trimmed Iterative Closest Point

Assumptions:

1. 2 sets of 3D points: **data** set  $\mathcal{P} = \{\mathbf{p}_i\}_1^{N_p}$  and **model** set  $\mathcal{M} = \{\mathbf{m}_i\}_1^{N_m}$ . ( $N_p \neq N_m$ .) Points may be surface as well as volumetric measurements.
2. Minimum guaranteed rate of data points that can be paired is known\*: **minimum overlap**  $\xi$ . Number of data points that can be paired  $N_{po} = \xi N_p$ .
3. **Rough pre-registration**: max initial relative rotation  $30^\circ$ .
4. **Overlapping part is characteristic enough** to allow for unambiguous matching\*\*:
  - no high symmetry
  - no 'featureless' data

\* If  $\xi$  is unknown, it is set automatically: run TrICP several times, select best result.

\*\* Typical for most registration algorithms.

## Problem statement and notation

**Informal statement:** Find Euclidean transformation that brings an  $N_{po}$ -point subset of  $\mathcal{P}$  into best possible alignment with  $\mathcal{M}$ .

For rotation  $\mathbf{R}$  and translation  $\mathbf{t}$ , transformed points of  $\mathcal{P}$  are

$$\mathbf{p}_i(\mathbf{R}, \mathbf{t}) = \mathbf{R}\mathbf{p}_i + \mathbf{t}, \quad \mathcal{P}(\mathbf{R}, \mathbf{t}) = \{\mathbf{p}_i(\mathbf{R}, \mathbf{t})\}_1^{N_p}$$

**Individual distance** from data point  $\mathbf{p}_i(\mathbf{R}, \mathbf{t})$  to  $\mathcal{M}$ :

$$\begin{aligned} \mathbf{m}_{cl}(i, \mathbf{R}, \mathbf{t}) &\doteq \arg \min_{\mathbf{m} \in \mathcal{M}} \|\mathbf{m} - \mathbf{p}_i(\mathbf{R}, \mathbf{t})\| \\ d_i(\mathbf{R}, \mathbf{t}) &\doteq \|\mathbf{m}_{cl}(i, \mathbf{R}, \mathbf{t}) - \mathbf{p}_i(\mathbf{R}, \mathbf{t})\| \end{aligned}$$

**Formal statement:** Find rigid motion  $(\mathbf{R}, \mathbf{t})$  that minimises sum of least  $N_{po}$  square distances  $d_i^2(\mathbf{R}, \mathbf{t})$ .

Conventional ICP:  $\xi = 1$  and  $N_{po} = N_p$ .

TrICP: smooth transition to ICP as  $\xi \rightarrow 1$ .

Basic idea of TRiCP: Consistent use of LTS in **deterministic** way.

Start with previous  $S'_{TS} = \text{huge\_number}$ .

Iterate until any of **stopping conditions** is satisfied.

### *Algorithm 2: Trimmed Iterative Closest Point*

1. **Closest point**: For each point  $\mathbf{p}_i \in \mathcal{P}$ , find closest point in  $\mathcal{M}$  and compute  $d_i^2$ .
  2. **Trimmed Squares**: Sort  $d_i^2$ , select  $N_{po}$  least values and calculate their sum  $S_{TS}$ .
  3. **Convergence test**: If any of stopping conditions is satisfied, exit; otherwise, set  $S'_{TS} = S_{TS}$  and continue.
  4. **Motion calculation**: For  $N_{po}$  selected pairs, compute optimal motion  $(\mathbf{R}, \mathbf{t})$  that minimises  $S_{TS}$ .
  5. **Data set transformation**: Transform  $\mathcal{P}$  by  $(\mathbf{R}, \mathbf{t})$  and go to 1.
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# Stopping conditions

- Maximum allowed number of iterations  $N_{iter}$  has been reached, **or**
- Trimmed MSE is sufficiently small, **or**
- Change of Trimmed MSE is sufficiently small.

**Trimmed MSE**  $e$ : For sorted distances  $d_{s1} \leq d_{s2} \leq \dots \leq d_{sN_{po}} \leq \dots \leq d_{sN_p}$ ,

$$S_{TS} \doteq \sum_{si=s1}^{sN_{po}} d_{si}^2 \quad e \doteq \frac{S_{TS}}{N_{po}}$$

**Change of Trimmed MSE:**  $|S_{TS} - S'_{TS}|$

# Implementation details

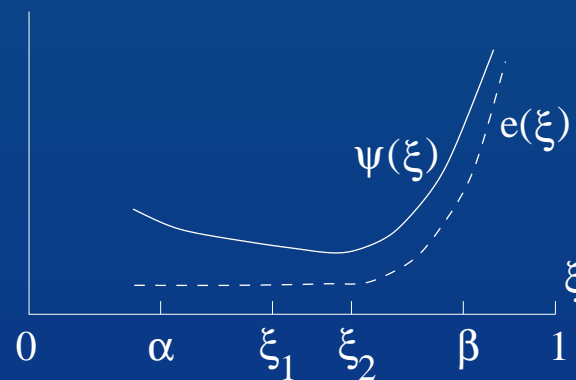
- Finding closest point: Use **boxing structure** (Chetverikov 1991) that partitions space into uniform boxes, cubes. Update box size as  $\mathcal{P}$  approaches  $\mathcal{M}$ .
  - simple
  - fast, especially at beginning of iterations
  - uses memory in inefficient way
- Sorting individual distances and calculating LTS: Use **heap sort**.
- Computing optimal motion: Use **Unit Quaternions**.
  - robust to noise
  - stable in presence of degenerate data ('flat' point sets)
  - relatively fast

## Automatic setting of overlap parameter $\xi$

When  $\xi$  is unknown, it is set automatically by minimising objective function

$$\psi(\xi) = \frac{e(\xi)}{\xi^{1+\lambda}}, \quad \lambda = 2$$

- $\psi(\xi)$  minimises trimmed MSE  $e(\xi)$  and tries to use as many points as possible.
- Larger  $\lambda$ : avoid undesirable alignments of symmetric and/or 'featureless' parts.
- $\psi(\xi)$  minimised using modified **Golden Section Search Algorithm**.



*Typical shapes of objective functions  $e(\xi)$  and  $\psi(\xi)$ .*

# Convergence

**Theorem:** *TrICP always converges monotonically to a local minimum with respect to trimmed MSE objective function.*

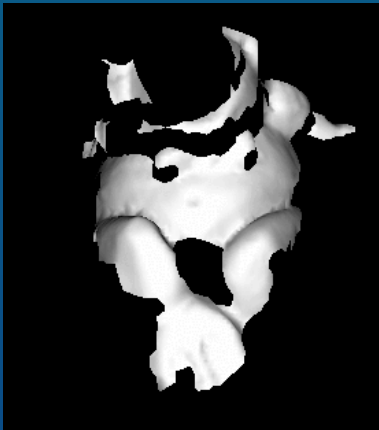
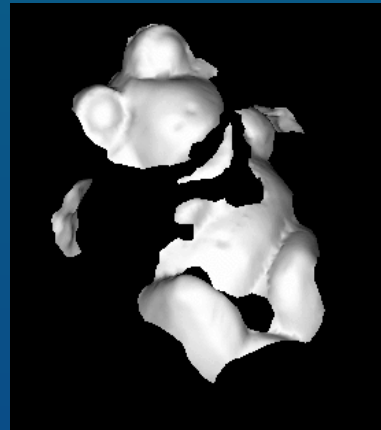
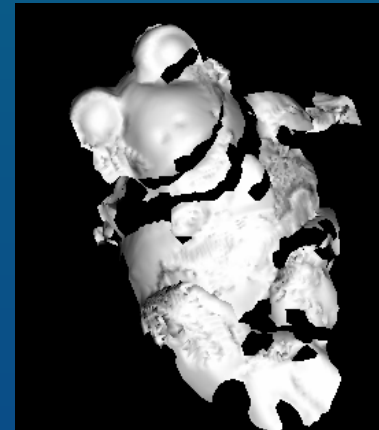
Sketch of proof:

- **Optimal motion** does not increase MSE: if it did, it would be inferior to identity transformation, as the latter does not change MSE.
- **Updating the closest points** does not increase MSE: no individual distance increases.
- **Updating the list of  $N_{po}$  least distances** does not increase MSE: to enter the list, any new pair has to substitute a pair with larger distance.
- Sequence of MSE values is nonincreasing and bounded below (by zero), hence it converges to a local minimum.

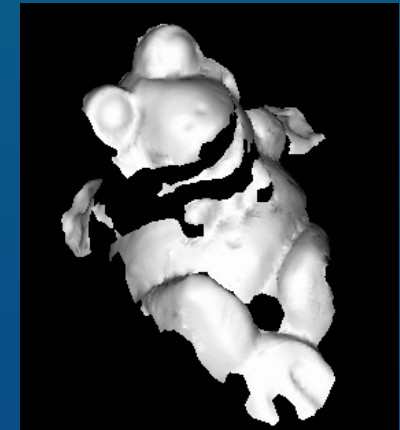
Convergence to global minimum depends on initial guess.



# Tests

set  $\mathcal{P}$ set  $\mathcal{M}$ 

result of ICP



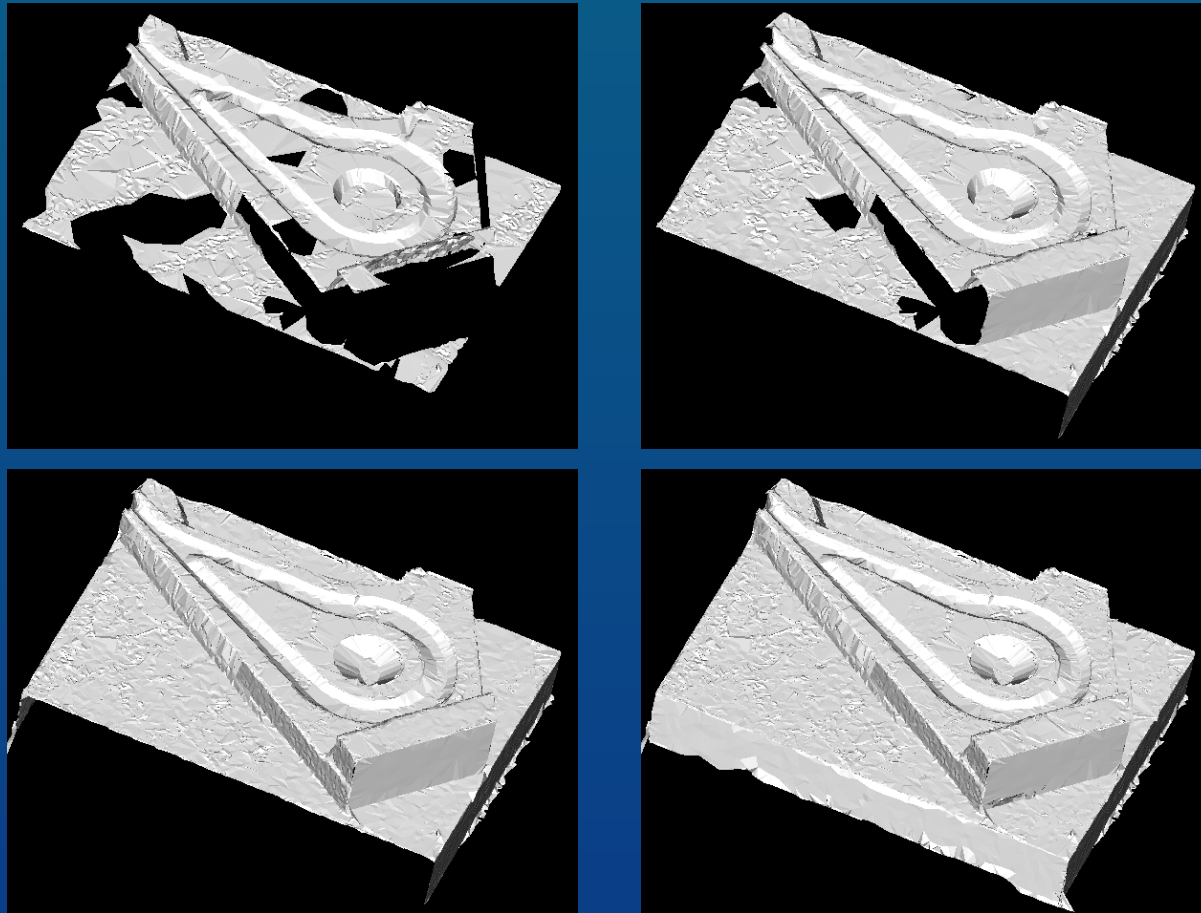
result of TrICP

Aligning *two partial measurements* of Frog. ( $\approx 3000$  points)  $\triangleleft \triangleleft$

*Numerical results for Frog data*

Method	$N_{iter}$	MSE	Exec.time*
ICP	45	5.83	7.4 sec
TrICP 70%	88	0.10	2.5 sec

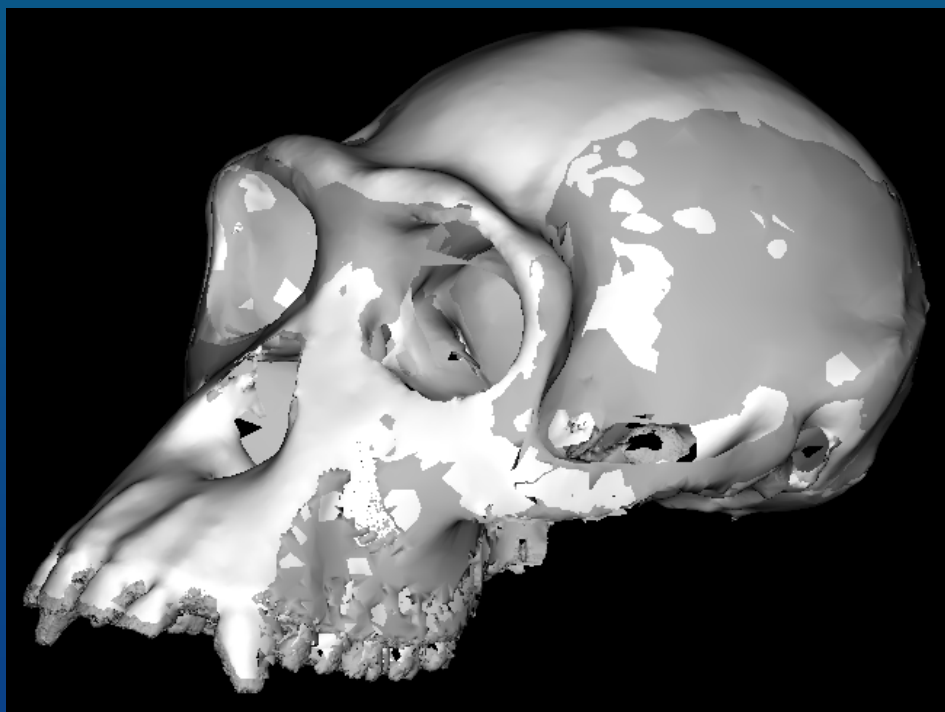
\* On 1.6 GHz PC



*Aligning four measurements of **Skoda part**. ( $\approx 6000$  points)*



*Aligning four measurements of **Fiat part**. (overlap  $\approx$  20%.)*

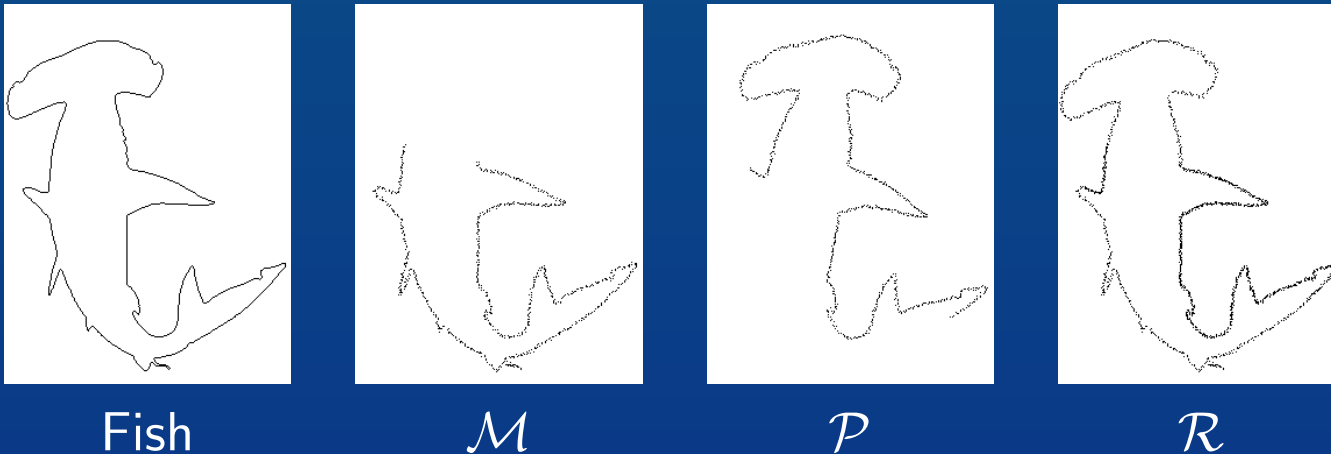


*Aligning two measurements of chimpanzee **Skull**. ( $\approx 100000$  points)*

# Comparing TrICP to ICRP for SQUID database

SQUID (University of Surrey, UK): 1100 shapes of different fishes

- $\mathcal{P}$  rotated by known angle. ( $1^\circ \dots 20^\circ$ )
- Different parts of  $\mathcal{M}$  and  $\mathcal{P}$  deleted.
- Noise added to both shapes.



*Aligning deteriorated **SQUID shapes**. Fish: original noise-free shape.*

*TrICP/ICRP errors for noisy SQUID data, degrees*

	100%	90%	80%	70%	60%
1°	0.05/0.05	0.08/0.06	0.07/0.08	0.10/0.12	0.19/0.23
5°	0.05/0.05	0.09/0.07	0.08/0.11	0.12/0.18	0.34/0.31
10°	0.05/0.05	0.09/0.07	0.10/0.18	0.19/0.60	0.58/1.70
15°	0.05/0.06	0.11/0.11	0.16/0.36	0.34/1.09	1.14/2.54
20°	0.05/0.11	0.10/0.16	0.20/0.49	0.69/1.51	1.79/3.03

- **ICRP** is efficient at small rotations and noise-free data.
- **TrICP** is more robust to rotations and incomplete, noisy data. Procedure for automatic setting of overlap available. Convergence proved.
- **Execution times** per alignment are comparable. Skoda 6000 pts: 8.3/7.2 sec, Skull 100000 pts: 38/182 sec.

# Future work

- Compare to other methods for a large set of shapes.
- To better avoid local minima, perturb initial orientation of data set.
- Extension to multiple point sets  $N > 2$ .
- Faster operation.
- More efficient usage of memory.